

# Coq-GT

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# Roadmap

Simplify

Contextual rewrite patterns

## simpl: my patch in a nutshell

Simpl unfolds a constant if it contains a fixpoint and the recursive argument evaluates to a constructor.

```
Eval simpl in (S x - y)
  match y with 0 => S x | S y' => x - y end
```

With my patch you can tell simpl to behave as follows:

```
Eval simpl in (S x - y)
  (S x - y)
Eval simpl in (S x - S y)
  (x - y)
Eval simpl in ((f \o g) x)
  f (g x)
Eval simpl in (f \o g)
  f \o g
```

In short, you choose how many arguments must be passed and which must evaluate to a constructor for a constant to be unfolded by simpl.

# Attaching meta-data to definition's arguments

Right now:

```
Definition fcomp A B C (f:B -> C) (g:A -> B) x := f (g x).
Infix "\o" := fcomp (at level 50, left associativity).
Implicit Arguments fcomp [A B C].
Arguments Scope fcomp [type_scope type_scope _ _ _ _].
Recursive Arguments fcomp [] 6.
```

What would be nice for the Math-Comp library

```
Arguments fcomp [[A%type B%type C%type]] f g x /.
```

Other examples:

```
Arguments nth [S%type] def !s%seq !i%nat.
Arguments foo [T%type] !n%nat !t / m g.
```

Because:

```
> grep '^Implicit Arguments' *.v | wc -l      #528
> grep '^Arguments Scope' *.v | wc -l        #158
```

# Why occurrence numbers are bad

```
...
g := [morphism of sdprodm defXA phiAiM] : {morphism
      joining_group A X >-> gT}
ker g : 'ker g = 'Mho^1(A)
skk : 'ker (coset ('ker g)) \subset 'ker g
nkA : joining_group A X \subset 'N('ker g)
fact_g := factm skk nkA : coset_groupType ('ker g) -> gT
imgX : X = fact_g @* (X / 'ker g)
nAA1 : A \subset 'N('Mho^1(A))
nXA1 : X \subset 'N('Mho^1(A))
=====
minnormal (fact_g @* (A / 'ker g)) X ->
minnormal (A / 'ker g) (X / 'ker g)

rewrite {1}imgX
```

# Why occurrence numbers are bad

```
...
g := [morphism of sdprodm defXA phiAiM] : {morphism
      joining_group A X >-> gT}
ker g : 'ker g = 'Mho^1(A)
skk : 'ker (coset ('ker g)) \subset 'ker g
nkA : joining_group A X \subset 'N('ker g)
fact_g := factm skk nkA : coset_groupType ('ker g) -> gT
imgX : X = fact_g @* (X / 'ker g)
nAA1 : A \subset 'N('Mho^1(A))
nXA1 : X \subset 'N('Mho^1(A))
=====
minnormal (fact_g @* (A / 'ker g)) X ->
minnormal (A / 'ker g) (X / 'ker g)
```

```
rewrite {1}imgX
```

```
rewrite {29}imgX
```



# Why occurrence numbers are bad

Occurrence numbers are bad for the following reasons:

- ▶ can be hard to write
- ▶ scripts are less informative when they break

SSR 1.3 contextual patterns:

- ▶ specify the occurrences looking at their context

```
minnormal (fact_g @* (A / 'ker g)) X ->  
  minnormal (A / 'ker g) (X / 'ker g)
```

One gives the context

```
rewrite [      minnormal (_ @* _) _]imgX
```



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SSR 1.3 contextual patterns:

- ▶ specify the occurrences looking at their context

```
minnormal (fact_g @* (A / 'ker g)) X ->  
  minnormal (A / 'ker g) (X / 'ker g)
```

One gives the context and names a hole

```
rewrite [R in minnormal (_ @* _) R] imgX
```

# Rewrite (contextual) patterns

## Terminology

`matching` head constant driven

`addnC` : `forall` a b, a + b = b + a

`redex` the term being rewritten, identified with a given pattern or a pattern inferred looking at the rule

## Rewrite pattern syntax

```
rewrite rule
```

```
rewrite [t]rule
```

```
rewrite [in t]rule
```

```
rewrite [X in t]rule
```

```
rewrite [in X in t]rule
```

```
rewrite [e in X in t]rule
```

```
rewrite [e as X in t]rule
```

# Rewrite patterns — example 1

The rule

`addnC` : `_ + _ = _ + _`

The tactic invocation

`rewrite addnC.`

The goal

`(x + y) + f x (x + y).+1 = 0`

## Rewrite patterns — example 2

The rule

$(\text{addnC } x.+1) : x.+1 + \_ = \_ + x.+1$

The tactic invocation

`rewrite [_.+1] (addnC x.+1).`

The goal

$(x + y) + f\ x \ \underline{(x + y).+1} = 0$

Because  $(x + y).+1 = x.+1 + \_$

## Contextual rewrite patterns — example 3

The rule

`addnC` : `_ + _ = _ + _`

The tactic invocation

`rewrite [in f _ _] addnC.`

The goal

`(x + y) + f x (x + y).+1 = 0`

## Contextual rewrite patterns — example 4

The rule

$$(\text{addnC } x.+1) : x.+1 + \_ = \_ + x.+1$$

The tactic invocation

```
rewrite [R in f _ R] (addnC x.+1).
```

The goal

$$(x + y) + f (x.+1 + y) \underline{(x + y).+1} = 0$$

Because R captured  $(x + y).+1 = x.+1 + \_$

## Contextual rewrite patterns — example 5

The rule

$(\text{addnC } x) : x + \_ = \_ + x$

The tactic invocation

`rewrite [in R in f _ R] (addnC x).`

The goal

$(x + y) + f\ x\ (z + \underline{x + y}).+1) = 0$

Because R captured  $z + (x + y).+1$

## Contextual rewrite patterns — example 6

The rule

$$(\text{addnC } x.+1) : x.+1 + \_ = \_ + x.+1$$

The tactic invocation

```
rewrite [_.+1 in R in f _ R] (addnC x.+1).
```

The goal

$$(x + y) + f\ x\ (z + \underline{(x + y).+1}) = 0$$

Because R captured  $z + (x + y).+1$  and  $_.+1$  matched

$$(x + y).+1 = x.+1 + \_$$



## Contextual rewrite patterns — example 7

The rule

`addnC : _ + _ = _ + _`

The tactic invocation

`rewrite [x.+1 + y as R in f _ (_ + R)]addnC.`

The goal

$(x + y) + f\ x\ (z + \underline{(x + y).+1}) = 0$

Because R captured  $(x + y).+1 = x.+1 + y = \_ + \_$

# Contextual patterns — the future

Patterns everywhere:

```
set t := {3}(a + _).  
set t := (a + _ in R in _ = R).
```

With user defined shortcuts

```
Notation RHS := (X in _ = X)%pattern.  
set t := (a + _ in RHS).  
move: (a + _ in RHS).  
rewrite {2}[in RHS]addnC.  
elim: (n in RHS).
```

Questions:

- ▶ Do you like this idea?
- ▶ Could this feature be part of Coq?

I've a branch in which the ssreflect plugin is split into an ssrmatching plugin and the ssreflect plugin builds on top of it.