Mathematical Knowledge Management in Algebraic Topology

Jónathan Heras Vicente

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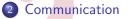
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4 Deduction

5 Conclusions and Further work



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Context

Mathematical Knowledge Management (MKM)

- Computation
- Deduction
- Communication



Context

Mathematical Knowledge Management (MKM)

- Computation
- Deduction
- Communication
- Algebraic Topology
 - Mathematical subject which studies topological spaces by algebraic means
 - Applications in Coding theory, Robotics, Digital Image analysis

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Context

Mathematical Knowledge Management (MKM)

- Computation
- Deduction
- Communication
- Algebraic Topology
 - Mathematical subject which studies topological spaces by algebraic means
 - Applications in Coding theory, Robotics, Digital Image analysis
- Kenzo
 - Computer Algebra system devoted to Algebraic Topology developed by F. Sergeraert
 - Common Lisp package
 - Homology groups unreachable by any other means

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- Communication
 - Human beings
 - Other programs



- Communication
 - Human beings
 - Other programs
- Computation
 - Kenzo



- Communication
 - Human beings
 - Other programs
- Computation
 - Kenzo
 - GAP



- Communication
 - Human beings
 - Other programs
- Computation
 - Kenzo
 - GAP
 - New modules



- Communication
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- Deduction



- Communication
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 - Kenzo
 - GAP
 - New modules
- Deduction
 - Certification of Kenzo algorithms
 - Isabelle/Hol
 - Coq



Communication

- Human beings
- Other programs
- Computation
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 - Certification of Kenzo algorithms
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 - Coq
 - Certification of Kenzo programs
 - ACL2

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5 Conclusions and Further work



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Motivation

Motivation

Algebraic Topology Expert Common Lisp Expert $\}$ \Rightarrow Makes the most of Kenzo



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Motivation

Motivation

Algebraic Topology Expert $\left.\right\} \Rightarrow$ Makes the most of Kenzo Common Lisp Expert

Non Common Lisp Expert \Rightarrow Unfriendly front-end

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Motivation

Algebraic Topology Expert Common Lisp Expert $\}$ \Rightarrow Makes the most of Kenzo

Non Common Lisp Expert \Rightarrow Unfriendly front-end

Non Algebraic Topology Expert \Rightarrow Needs guidance

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• Develop a system which increases Kenzo accessibility



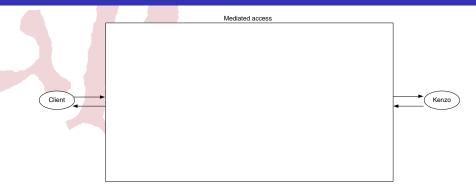
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- Develop a system which increases Kenzo accessibility
 - Friendly front-end
 - Mediated access to Kenzo



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Representation of mathematical knowledge

 $\left. \begin{array}{l} \mbox{Independent of programming language} \\ \mbox{Easily interchangeable} \end{array} \right\} \Rightarrow XML \Rightarrow \left\{ \begin{array}{l} \mbox{OpenMath} \\ \mbox{MathML} \\ \mbox{New language} \end{array} \right.$

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XML-Kenzo

- construction and computation requests
- Includes some mathematical knowledge
 - Small Type System: CC, SS, SG, ASG
 - Some restrictions about arguments (Sⁿ)

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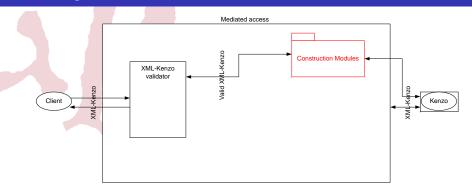
XML-Kenzo validator

- Common Lisp module
- ${lackstyle}$ Receives XML-Kenzo requests \rightarrow validates them against XML schema definition
 - Type restrictions
 - Mathematical restrictions
 - Kenzo restrictions

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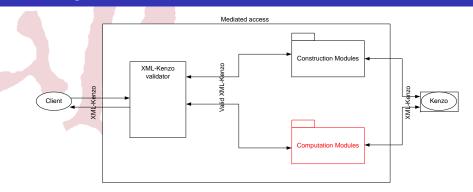
Construction modules

- Process construction requests
- Check restrictions not included in XML-Kenzo (functional dependencies)
- 4 modules: CC, SS, SG and ASG

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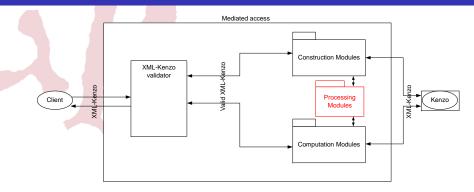
Computation modules

- Process computation requests
- Check restrictions not included in XML-Kenzo (reduction degree)
- 2 modules: Homology and Homotopy

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Processing modules

- Help construction and computation modules
- 2 modules: Reduction degree and Homotopy assistant

The whole framework

Towards the whole framework



- Mediated access to Kenzo
- Friendly front-end
- but . . .



Towards the whole framework

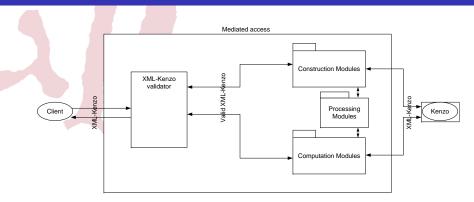
This architecture provides

- Mediated access to Kenzo
- Friendly front-end
- but . . .

Desirable features

- Different clients
- Efficient
- Extensible
- Adaptable to different needs

Towards different clients

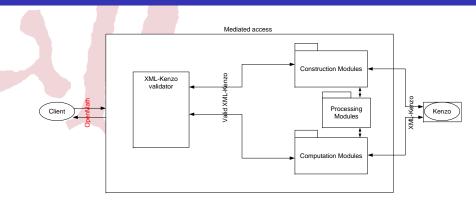


• XML-Kenzo is ad-hoc



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Towards different clients



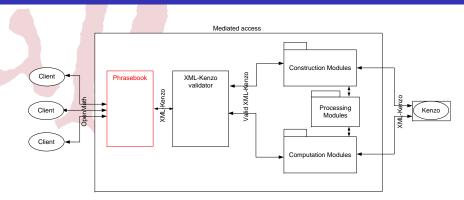
- XML-Kenzo is ad-hoc
- OpenMath
 - Standard
 - Communication with the outside world

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Towards different clients



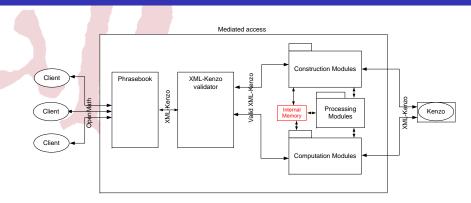
- XML-Kenzo is ad-hoc
- OpenMath
 - Standard
 - Communication with the outside world
- Phrasebook

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Towards Efficiency



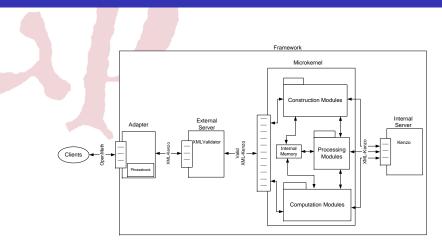
- System roughly equivalent to Kenzo
- Internal memory
 - Memoization technique
 - Store spaces and computations

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The Kenzo framework



Based on well-known patterns and methods

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Towards extensibility

Include new functionality



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Towards extensibility

Include new functionality

Increase Kenzo functionality



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Towards extensibility

Include new functionality

- Increase Kenzo functionality
- Include new systems
 - Computer Algebra systems
 - Theorem Prover tools



Towards extensibility

Include new functionality

- Increase Kenzo functionality
- Include new systems
 - Computer Algebra systems
 - Theorem Prover tools
- Interoperability

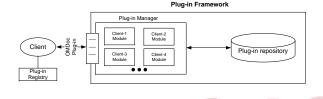


Towards extensibility

Include new functionality

- Increase Kenzo functionality
- Include new systems
 - Computer Algebra systems
 - Theorem Prover tools
- Interoperability

Kenzo framework as a client of a plug-in framework



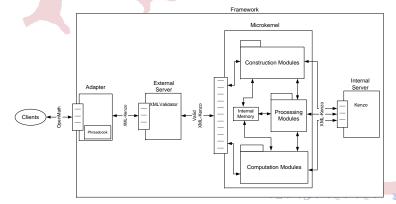
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Including new Kenzo functionality

```
<code id="new-constructor">
   <data format="Kf/internal-server"> new-constructor-is.lisp </data>
   <data format="Kf/microkernel"> new-constructor-m.lisp </data>
   <data format="Kf/external-server"> XML-Kenzo.xsd </data>
   <data format="Kf/adapter"> new-constructor-a.lisp </data>
   <data format="Kf/adapter"> new-constructor-a.lisp </data>
   <data format="Kf/adapter"> new-constructor-a.lisp </data>
   <data format="Kf/adapter"> new-constructor-a.lisp </data>
   </data>
   <data format="Kf/adapter"> new-constructor-a.lisp </data>
   </data>
   </data</pre>
```



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• A Graphical User Interface implemented in Common Lisp



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- A Graphical User Interface implemented in Common Lisp
- Design decisions
 - Functionality (Common Lisp) + Structure (XUL) + Layout (stylesheet)



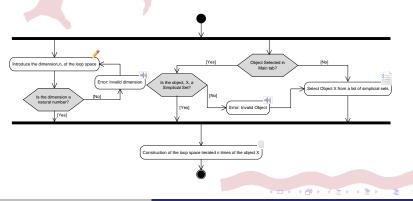
- A Graphical User Interface implemented in Common Lisp
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 - Guided by heuristics



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 - Design of interaction: Noesis method



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A customizable GUI

- Features
 - Extensibility
 - Adaptability



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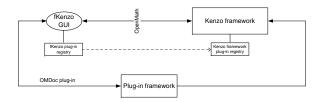
A customizable GUI

- Features
 - Extensibility
 - Adaptability
- The Graphical User Interface is
 - client of Kenzo framework
 - client of plug-in framework
 - GUI organized through modules: basic and experimental



A customizable GUI

- Features
 - Extensibility
 - Adaptability
- The Graphical User Interface is
 - client of Kenzo framework
 - client of plug-in framework
 - GUI organized through modules: basic and experimental



The whole system is called *fKenzo*

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fKenzo

fKenzo: A tool adaptable to different needs

• Beginner student of Algebraic Topology

Load and remove basic modules



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fKenzo: A tool adaptable to different needs

- Beginner student of Algebraic Topology
 - Load and remove basic modules
- Advanced student or researcher of Algebraic Topology
 - Load and remove basic modules
 - Load and remove experimental modules



fKenzo: A tool adaptable to different needs

- Beginner student of Algebraic Topology
 - Load and remove basic modules
- Advanced student or researcher of Algebraic Topology
 - Load and remove basic modules
 - Load and remove experimental modules
 - Digital images
 - GAP
 - ACL2
 - . . .







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fKenzo

fKenzo summary

- fKenzo
 - Integral assistant for Algebraic Topology



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- fKenzo
 - Integral assistant for Algebraic Topology
 - Constrains Kenzo functionality



- fKenzo
 - Integral assistant for Algebraic Topology
 - Constrains Kenzo functionality
 - Provides guidance in the interaction and a friendly front-end



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 - Extensible



- fKenzo
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 - Provides guidance in the interaction and a friendly front-end
 - Extensible
 - Adaptable to different needs

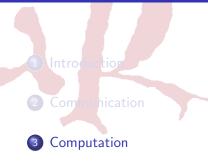


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- fKenzo
 - Integral assistant for Algebraic Topology
 - Constrains Kenzo functionality
 - Provides guidance in the interaction and a friendly front-end
 - Extensible
 - Adaptable to different needs
- Our architecture can be applied in other contexts



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4 Deduction

5 Conclusions and Further work



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- Increase Kenzo computational capabilities
 - Simplicial Complexes
 - Digital Images
 - Pushout of Simplicial Sets



- Increase Kenzo computational capabilities
 - Simplicial Complexes
 - Digital Images
 - Pushout of Simplicial Sets
- Integrated in *fKenzo*

- Increase Kenzo computational capabilities
 - Simplicial Complexes
 - Digital Images
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- Integrated in *fKenzo*
- Verified using ACL2



- Increase Kenzo computational capabilities
 - Simplicial Complexes
 - Digital Images
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- Integrated in *fKenzo*
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Compute homology groups of a space X

```
• H_*(X) = H_*(C_*(X))
```



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Compute homology groups of a space X

- $H_*(X) = H_*(C_*(X))$
 - C_{*}(X) has finite nature
 - Differential maps can be expressed as integer matrices
 - Homology groups: Smith Normal Form



Compute homology groups of a space X

- $H_*(X) = H_*(C_*(X))$
 - C_{*}(X) has finite nature
 - Differential maps can be expressed as integer matrices
 - Homology groups: Smith Normal Form
 - C_{*}(X) has non finite nature
 - Previous methods cannot be applied
 - Effective Homology
 - · Sergeraert's ideas
 - \cdot Provides real algorithms to compute homology groups
 - · Implemented in the Kenzo system

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Compute homology groups of a space X

- $H_*(X) = H_*(C_*(X))$
 - C_{*}(X) has finite nature (Effective Chain Complex)
 - Differential maps can be expressed as integer matrices
 - Homology groups: Smith Normal Form
 - C_{*}(X) has non finite nature (Locally Effective Chain Complex)
 - Previous methods cannot be applied
 - Effective Homology
 - \cdot Sergeraert's ideas
 - \cdot Provides real algorithms to compute homology groups
 - · Implemented in the Kenzo system

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Definition

An effective chain complex is a free chain complex of \mathbb{Z} -modules, $C_* = (C_n, d_n)_{n \in \mathbb{N}}$, where each group C_n is finitely generated and

- an algorithm returns a Z-base in each grade n
- an algorithm provides the differentials d_n

Definition

A locally effective chain complex is a free chain complex of \mathbb{Z} -modules $C_* = (C_n, d_n)_{n \in \mathbb{N}}$ where each group C_n can have infinite nature, but there exists an algorithm such that $\forall x \in C_n$, we can compute $d_n(x)$

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Definition

An effective chain complex is a free chain complex of \mathbb{Z} -modules, $C_* = (C_n, d_n)_{n \in \mathbb{N}}$, where each group C_n is finitely generated and

- an algorithm returns a Z-base in each grade n
- an algorithm provides the differentials d_n
- differentials $d_n : C_n \to C_{n-1}^r$ can be expressed as integer matrices
- possible to compute Ker d_n and Im d_{n+1}
- possible to compute the homology groups

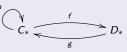
Definition

A locally effective chain complex is a free chain complex of \mathbb{Z} -modules $C_* = (C_n, d_n)_{n \in \mathbb{N}}$ where each group C_n can have infinite nature, but there exists an algorithm such that $\forall x \in C_n$, we can compute $d_n(x)$

- impossible to compute Ker d_n and Im d_{n+1}
- possible to perform local computations, differential of a generator

Definition

A reduction ρ between two chain complexes C_* y D_* (denoted by $\rho : C_* \Rightarrow D_*$) is a triple $\rho = (f, g, h)$



satisfying the following relations

1) $fg = Id_{D_*}$;

2)
$$d_C h + h d_C = \operatorname{Id}_{C_*} - gf;$$

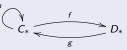
3)
$$fh = 0;$$
 $hg = 0;$ $hh = 0.$



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Definition

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3)
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 $hg = 0;$ $hh = 0.$

Theorem

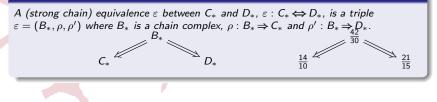
If $C_* \Rightarrow D_*$, then $C_* \cong D_* \oplus A_*$, with A_* acyclic, which implies that $H_n(C_*) \cong H_n(D_*)$ for all n.

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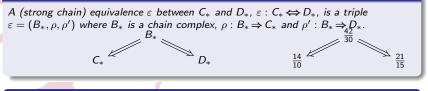
Definition





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Definition



Definition

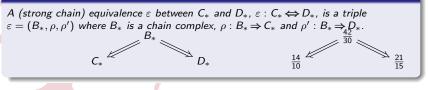
An object with effective homology is a quadruple $(X, C_*(X), HC_*, \varepsilon)$ where

- X is a locally effective object
- C_{*}(X) is a (locally effective) chain complex canonically associated with X, which allows the study of the homological nature of X
- HC_{*} is an effective chain complex
- ε is a equivalence ε : $C_*(X) \Leftrightarrow HC_*$

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Effective Homology preliminaries

Definition



Definition

An object with effective homology is a quadruple $(X, C_*(X), HC_*, \varepsilon)$ where

- X is a locally effective object
- C_{*}(X) is a (locally effective) chain complex canonically associated with X, which allows the study of the homological nature of X
- HC_{*} is an effective chain complex
- ε is a equivalence $\varepsilon : C_*(X) \Leftrightarrow HC_*$

Theorem

Let an object with effective homology $(X, C_*(X), HC_*, \varepsilon)$ then $H_n(X) \cong H_n(HC_*)$ for all n.

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Definition

Let f, g morphisms,





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Definition

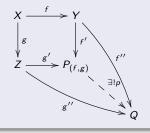
Let f, g morphisms, the pushout of f, g is an object $P_{(f,g)}$ for which the diagram





Definition

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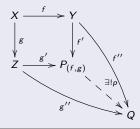


- commutes
- respects the universal property



Definition

Let f, g morphisms, the pushout of f, g is an object $P_{(f,g)}$ for which the diagram



- commutes
- respects the universal property

Standard Construction

$$\mathcal{P}_{(f,g)}\cong (Y\amalg (X imes I)\amalg Z)/\sim$$
 where

I is the unit interval

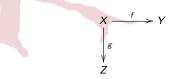
• for every
$$x \in X$$
, ~

•
$$(x, 0) \sim f(x) \in Y$$

•
$$(x,1) \sim g(x) \in Z$$

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Given $f: X \rightarrow Y$ and $g: X \rightarrow Z$ simplicial morphisms where X, Y and Z are simplicial sets





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Given $f: X \rightarrow Y$ and $g: X \rightarrow Z$ simplicial morphisms where X, Y and Z are simplicial sets



Given $f: X \rightarrow Y$ and $g: X \rightarrow Z$ simplicial morphisms where X, Y and Z are simplicial sets

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow g & & \downarrow f' \\ Z & \xrightarrow{g'} & P_{(f,g)} \end{array}$$

Algorithm (Standard Construction)

Input: two simplicial morphisms $f: X \to Y$ and $g: X \to Z$ where X, Y and Z are simplicial sets *Output*: the pushout $P_{(f,g)}$, a simplicial set

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Given $f: X \rightarrow Y$ and $g: X \rightarrow Z$ simplicial morphisms where X, Y and Z are simplicial sets with effective homology

$$(X, C_*(X), HX_*, \varepsilon_X) \xrightarrow{f} (Y, C_*(Y), HY_*, \varepsilon_Y)$$

$$\downarrow^g$$

$$(Z, C_*(Z), HZ_*, \varepsilon_Z)$$



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Given $f: X \to Y$ and $g: X \to Z$ simplicial morphisms where X, Y and Z are simplicial sets with effective homology



Given $f: X \to Y$ and $g: X \to Z$ simplicial morphisms where X, Y and Z are simplicial sets with effective homology



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Given $f: X \to Y$ and $g: X \to Z$ simplicial morphisms where X, Y and Z are simplicial sets with effective homology

Algorithm (joint work with F. Sergeraert)

Input: two simplicial morphisms $f: X \to Y$ and $g: X \to Z$ where X, Y and Z are simplicial sets with effective homology Output: the effective homology version of $P_{(f,g)}$, that is, an equivalence $C_*(P_{(f,g)}) \Leftrightarrow HP_*$, where HP_* is an effective chain complex

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Sketch of the algorithm

Theorem (SES Theorems)

Let

$$0 \stackrel{0}{\longleftrightarrow} A_* \stackrel{\sigma}{\underbrace{\longleftrightarrow}} B_* \stackrel{\rho}{\underbrace{\longleftarrow}} C_* \longleftarrow 0$$

be an effective short exact sequence of chain complexes. Then three general algorithms are available

 $\begin{array}{l} SES_1 : (B_{*,EH}, C_{*,EH}) \mapsto A_{*,EH} \\ SES_2 : (A_{*,EH}, C_{*,EH}) \mapsto B_{*,EH} \\ SES_3 : (A_{*,EH}, B_{*,EH}) \mapsto C_{*,EH} \end{array}$

producing the effective homology of one chain complex when the effective homology of both others is given.



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Sketch of the algorithm

Theorem (SES Theorems)

Let

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producing the effective homology of one chain complex when the effective homology of both others is given.

$$0 \longleftarrow M_* \xrightarrow[j]{\sigma} C_*P \xrightarrow{\rho} C_*Y \oplus C_*Z \longleftarrow 0$$

 M_* is the chain complex associated with $X \times \Delta^1$ but with the simplexes of $X \times (0)$ and $X \times (1)$ cancelled

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Sketch of the algorithm continued

- Step 1. From $f : X \to Y$ and $g : X \to Z$ simplicial morphisms, $P_{(f,g)}$ and its associated chain complex C_*P are constructed
- Step 2. The effective homology of *M*_{*} is constructed • Define

$$0 \longleftarrow M_* \xleftarrow{\sigma_2}{j_2} C_*(X \times \Delta^1) \xleftarrow{\rho_2}{j_2} C_*(X \times (0)) \oplus C_*(X \times (1)) \longleftarrow 0$$

- Construct the chain complex M_{*}
- Build the effective homology of $C_*(X \times \Delta^1)$
- Construct the effective homology of C_{*} (X × (0)) ⊕ C_{*} (X × (1))
- Construct the effective homology of Cone(i2)
- Construct the reduction M_{*} ⇐ Cone(i2)
- Obtain the effective homology of M_{*} applying case SES₁
- Step 3. The effective homology of $C_*X \oplus C_*Y$ is constructed
- Step 4. The effective homology of the pushout $P_{(f,g)}$ is constructed
 - Define

$$0 \longleftarrow M_* \xleftarrow{\sigma} C_* P \xleftarrow{\rho} C_* Y \oplus C_* Z \longleftarrow 0$$

- Construct the effective homology of $(C_* Y \oplus C_* Z)^{[1]}$
- Define the morphism shift : $C_* Y \oplus C_* Z \to (C_* Y \oplus C_* Z)^{[1]}$
- Define the chain complex morphism χ : M_{*} → (C_{*} Y ⊕ C_{*}Z)^[1]
- Construct the effective homology of Cone(χ)
- Obtain the effective homology of P_(f,g) applying case SES₂

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Mathematical Knowledge Management in Algebraic Topology

(同) (王) (王)

$SL_2(\mathbb{Z})$

• Group of 2 \times 2 matrices with determinant 1 over $\mathbb Z$



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$SL_2(\mathbb{Z})$

- Group of 2 \times 2 matrices with determinant 1 over $\mathbb Z$
- Isomorphic to $\mathbb{Z}_4 *_{\mathbb{Z}_2} \mathbb{Z}_6$
 - J. P. Serre. *Trees.* Springer-Verlag, 1980



$SL_2(\mathbb{Z})$

- Group of 2 \times 2 matrices with determinant 1 over $\mathbb Z$
- Isomorphic to $\mathbb{Z}_4 *_{\mathbb{Z}_2} \mathbb{Z}_6$

J. P. Serre. *Trees.* Springer-Verlag, 1980

•
$$\mathcal{K}(\mathbb{Z}_2, 1) \xrightarrow{i_1} \mathcal{K}(\mathbb{Z}_4, 1)$$

 $\downarrow^{i_2} \qquad \qquad \downarrow$
 $\mathcal{K}(\mathbb{Z}_6, 1) \longrightarrow \mathcal{K}(SL_2(\mathbb{Z}), 1)$

🦫 K. S. Brown. Cohomology of Groups. Springer-Verlag, 1982

(同) (ヨ) (ヨ)

Construction of the object $K(SL_2(\mathbb{Z}, 1))$



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Construction of the object $K(SL_2(\mathbb{Z}, 1))$

Firstly, we define K(Z₂, 1), K(Z₄, 1) and K(Z₆, 1)
> (setf kz2 (k-zp-1 2)) ★
[K2 Abelian-Simplicial-Group]
> (setf kz4 (k-zp-1 4)) ★
[K15 Abelian-Simplicial-Group]
> (setf kz6 (k-zp-1 6)) ★
[K28 Abelian-Simplicial-Group]



Construction of the object $K(SL_2(\mathbb{Z},1))$

• Firstly, we define $K(\mathbb{Z}_2, 1), K(\mathbb{Z}_4, 1)$ and $K(\mathbb{Z}_6, 1)$

> (setf kz2 (k-zp-1 2)) \ [K2 Abelian-Simplicial-Group] > (setf kz4 (k-zp-1 4)) \ [K15 Abelian-Simplicial-Group] > (setf kz6 (k-zp-1 6)) \ [K28 Abelian-Simplicial-Group]

• Subsequently, we define the two simplicial morphisms i_1 and i_2

> (setf i1 (kzps-incl 2 4)) [K40 Simplicial-Morphism K2 -> K15] > (setf i2 (kzps-incl 2 6)) [K41 Simplicial-Morphism K2 -> K28]

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Construction of the object $K(SL_2(\mathbb{Z},1))$

• Firstly, we define $K(\mathbb{Z}_2, 1), K(\mathbb{Z}_4, 1)$ and $K(\mathbb{Z}_6, 1)$

> (setf kz2 (k-zp-1 2)) \ [K2 Abelian-Simplicial-Group] > (setf kz4 (k-zp-1 4)) \ [K15 Abelian-Simplicial-Group] > (setf kz6 (k-zp-1 6)) \ [K28 Abelian-Simplicial-Group]

- Subsequently, we define the two simplicial morphisms i₁ and i₂
 > (setf i1 (kzps-incl 2 4)) ★
 [K40 Simplicial-Morphism K2 -> K15]
 > (setf i2 (kzps-incl 2 6)) ★
 [K41 Simplicial-Morphism K2 -> K28]
- Finally, we construct the pushout of i₁ and i₂ (K(SL₂(ℤ), 1))
 > (setf ksl2z (pushout i1 i2)) ႃ
 [K52 Simplicial-Set]

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Computation of $\pi_*(\Sigma(SL_2(\mathbb{Z})))$



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Computation of $\pi_*(\Sigma(SL_2(\mathbb{Z})))$

- Firstly, we define $\Sigma(K(SL_2(\mathbb{Z}), 1))$
 - > (setf sksl2z (suspension ksl2z)) 🗜 [K62 Simplicial-Set]



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Computation of $\pi_*(\Sigma(SL_2(\mathbb{Z})))$

- Firstly, we define $\Sigma(K(SL_2(\mathbb{Z}), 1))$
 - > (setf sksl2z (suspension ksl2z)) 🗜 [K62 Simplicial-Set]
- Computing Homotopy groups (Hurewicz theorem)
 - > (homology sksl2z 1 3) 🕏 Homology in dimension 1:

Homology in dimension 2: Component Z/12Z



Computation of $\pi_*(\Sigma(SL_2(\mathbb{Z})))$

- Firstly, we define $\Sigma(K(SL_2(\mathbb{Z}), 1))$
 - > (setf sksl2z (suspension ksl2z)) 🗜 [K62 Simplicial-Set]
- Computing Homotopy groups (Hurewicz theorem)
 - > (homology sksl2z 1 3) Homology in dimension 1:

Homology in dimension 2: Component Z/12Z

•
$$\pi_1(\Sigma(SL_2(\mathbb{Z}))) = 0, \ \pi_2(\Sigma(SL_2(\mathbb{Z}))) = \mathbb{Z}/12\mathbb{Z}$$

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```
    Computing Homotopy groups continued (Whitehead tower)

  > (setf tau (zp-whitehead 12 sksl2z (chml-clss sksl2z 2))) 🕂
  [K91 Fibration K62 -> K79]
  > (setf x (fibration-total tau))
  [K97 Simplicial-Set]
  > (homology x 3)
  Homology in dimension 3:
  Component Z/12Z
  > (setf tau2 (zp-whitehead 12 x (chml-clss x 3))) 🗜
  [K228 Fibration K97 -> K214]
  > (setf x2 (fibration-total tau2)) 🖌
  [K234 Simplicial-Set]
  > (homology x2 4)
  Homology in dimension 4:
  Component Z/12Z
  Component Z/2Z
  > ...
```

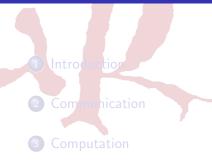
```
    Computing Homotopy groups continued (Whitehead tower)

  > (setf tau (zp-whitehead 12 sksl2z (chml-clss sksl2z 2))) 🕂
  [K91 Fibration K62 -> K79]
  > (setf x (fibration-total tau)) 🕂
  [K97 Simplicial-Set]
  > (homology x 3)
  Homology in dimension 3:
  Component Z/12Z
  > (setf tau2 (zp-whitehead 12 x (chml-clss x 3))) 😾
  [K228 Fibration K97 -> K214]
  > (setf x2 (fibration-total tau2)) 🖁
  [K234 Simplicial-Set]
  > (homology x2 4)
  Homology in dimension 4:
  Component Z/12Z
  Component Z/2Z
  > ...
```

```
• \pi_3(\Sigma(SL_2(\mathbb{Z}))) = \mathbb{Z}/12\mathbb{Z}, \ \pi_4(\Sigma(SL_2(\mathbb{Z}))) = \mathbb{Z}/12\mathbb{Z} \oplus \mathbb{Z}/2\mathbb{Z}, \ldots
```

Deduction

Table of Contents



4 Deduction

5 Conclusions and Further work



Jónathan Heras Vicente



• ACL2 (A Computational Logic for an Applicative Common Lisp)



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ACL2

• ACL2 (A Computational Logic for an Applicative Common Lisp)

- Programming Language
- First-Order Logic
- Theorem Prover



ACL2

• ACL2 (A Computational Logic for an Applicative Common Lisp)

- Programming Language
- First-Order Logic
- Theorem Prover
- Proof techniques
 - Simplification
 - Induction
 - "The Method"



ACL2

• ACL2 (A Computational Logic for an Applicative Common Lisp)

- Programming Language
- First-Order Logic
- Theorem Prover
- Proof techniques
 - Simplification
 - Induction
 - "The Method"
- Encapsulate principle
 - Simulation of Higher-Order Logic



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Goal

Goal

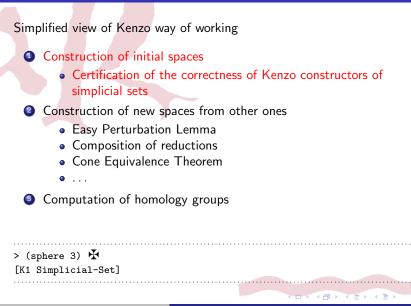
Simplified view of Kenzo way of working

- Construction of initial spaces
- Construction of new spaces from other ones
- Computation of homology groups



Goal

Goal



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Mathematical context: Simplicial Sets

Definition

A simplicial set K, is a union $K = \bigcup_{q \ge 0} K^q$, where the K^q are disjoints sets, together with functions

$$\begin{array}{ll} \partial_i^q: \mathsf{K}^q \to \mathsf{K}^{q-1}, & q > 0, & i = 0, \dots, q, \\ \eta_i^q: \mathsf{K}^q \to \mathsf{K}^{q+1}, & q \ge 0, & i = 0, \dots, q, \end{array}$$

subject to the relations

$$\begin{array}{rcl} (1) & \partial_{i}^{q-1}\partial_{j}^{q} & = & \partial_{j-1}^{q-1}\partial_{i}^{q} & \mbox{if} & i < j, \\ (2) & \eta_{i}^{q+1}\eta_{j}^{q} & = & \eta_{j+1}^{q+1}\eta_{i}^{q} & \mbox{if} & i \leq j, \\ (3) & \partial_{i}^{q+1}\eta_{j}^{q} & = & \eta_{j-1}^{q-1}\partial_{i}^{q} & \mbox{if} & i < j, \\ (4) & \partial_{i}^{q+1}\eta_{i}^{q} & = & \mbox{identity} & = & \partial_{i+1}^{q+1}\eta_{i}^{q}, \\ (5) & \partial_{i}^{q+1}\eta_{j}^{q} & = & \eta_{j}^{q-1}\partial_{i-1}^{q} & \mbox{if} & i > j+1, \\ \end{array}$$



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The elements of K^q are called q-simplexes

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Mathematical Knowledge Management in Algebraic Topology

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Mathematical context: Simplicial Sets

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subject to the relations

- The elements of K^q are called q-simplexes
- A q-simplex x is degenerate if $x = \eta_i^{q-1}y$ for some simplex $y \in K^{q-1}$

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Mathematical context: Simplicial Sets

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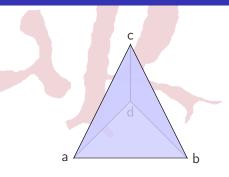
subject to the relations

- The elements of K^q are called q-simplexes
- A q-simplex x is degenerate if $x = \eta_i^{q-1}y$ for some simplex $y \in K^{q-1}$
- Otherwise x is called non-degenerate (or geometric)

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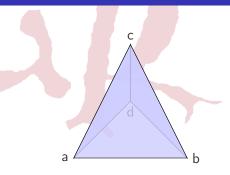
Mathematical context: Example



- 0-simplexes: vertices:
 (a), (b), (c), (d)
- non-degenerate 1-simplexes:
 edges:
 (a b),(a c),(a d),(b c),(b d),(c d)
- non-degenerate 2-simplexes: (filled) triangles: (a b c),(a b d),(a c d),(b c d)
- non-degenerate 3-simplexes:
 (filled) tetrahedron: (a b c d)

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Mathematical context: Example

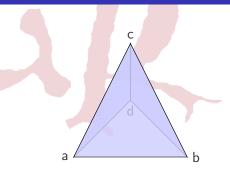


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face:
$$\partial_i(a \ b \ c) = \begin{cases} (b \ c) & \text{if } i = 0\\ (a \ c) & \text{if } i = 1\\ (a \ b) & \text{if } i = 2 \end{cases}$$

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Mathematical context: Example



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degeneracy:
$$\eta_i(a \ b \ c) = \begin{cases} (a \ a \ b \ c) & \text{if } i = 0 \\ (a \ b \ b \ c) & \text{if } i = 1 \\ (a \ b \ c \ c) & \text{if } i = 2 \end{cases}$$
 non-geometrical meaning

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Mathematical context: abstract simplexes

Proposition

Let K be a simplicial set. Any n-simplex $x \in K^n$ can be expressed in a unique way as a (possibly) iterated degeneracy of a non-degenerate simplex y in the following way

$$x = \eta_{j_k} \dots \eta_{j_1} y$$

with $y \in K^r$, $k = n - r \ge 0$, and $0 \le j_1 < \cdots < j_k < n$.



Mathematical context: abstract simplexes

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abstract simplex



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Mathematical context: abstract simplexes

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abstract simplex

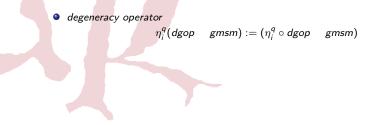
- Examples

$$\begin{array}{rll} \text{simplex} & \text{abstract simplex} \\ \text{non-degenerate} & (a \ b) & (\emptyset \ (a \ b)) \\ \text{degenerate} & (a \ a \ b \ c) & (\eta_0 \ (a \ b \ c)) \end{array}$$

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degeneracy operator

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$$q_i^q(dgop gmsm) := (\eta_i^q \circ dgop gmsm)$$

face operator

where

$$r = q - \{$$
number of degeneracies in *dgop* $\}$ and
 $k = i - \{$ number of degeneracies in *dgop* with index lower than $i\}$

degeneracy operator

$$q^q(dgop gmsm) := (\eta^q_i \circ dgop gmsm)$$

• face operator

$$\partial_i^q(\textit{dgop} \quad \textit{gmsm}) := \left\{ \begin{array}{ll} (\partial_i^q \circ \textit{dgop} \quad \textit{gmsm}) & \text{if} \quad \eta_i \in \textit{dgop} \lor \eta_{i-1} \in \textit{dgop} \\ (\partial_i^q \circ \textit{dgop} \quad \partial_k^r \textit{gmsm}) & \text{otherwise}; \end{array} \right.$$

where

 $r = q - \{$ number of degeneracies in $dgop \}$ and

 $k = i - \{$ number of degeneracies in *dgop* with index lower than $i\}$

invariant operator

 $(dgop gmsm) \in K^q$

- (length dgop) < q
- $gmsm \in K^r$ where r = q (length dgop)
- index of the first degeneracy in *dgop* is lower than *q*

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degeneracy operator

 $\eta_i^q(dgop \quad gmsm) := (\eta_i^q \circ dgop \quad gmsm)$

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 $\partial_i^q (\textit{dgop} \ \textit{gmsm}) := \left\{ \begin{array}{cc} (\partial_i^q \circ \textit{dgop} \ \textit{gmsm}) & \text{if} \ \eta_i \in \textit{dgop} \lor \eta_{i-1} \in \textit{dgop} \\ (\partial_i^q \circ \textit{dgop} \ \partial_k^r \textit{gmsm}) & \text{otherwise}; \end{array} \right.$

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- (length dgop) < q
- $gmsm \in K^r$ where r = q (length dgop)
- index of the first degeneracy in *dgop* is lower than *q*

Dependent part from the chosen simplicial set \rightarrow Affect geometric part

Independent parts from the chosen simplicial set \rightarrow Not affect geometric part

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- Functions
 - face-absm
 - degeneracy-absm
 - invariant-absm



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- Functions
 - face-absm
 - degeneracy-absm
 - invariant-absm

Dependent parts

- face-gmsm
- invariant-gmsm



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- Functions
 - face-absm
 - degeneracy-absm
 - invariant-absm

Dependent parts

- face-gmsm
- invariant-gmsm

Independent parts

- degeneracy
- face-independent
- invariant-independent



- Functions
 - face-absm
 - degeneracy-absm
 - invariant-absm
- Properties

$$\begin{array}{l} \bullet \quad \partial_{i}^{q-1}\partial_{j}^{q} = \partial_{j-1}^{q-1}\partial_{i}^{q} \text{ if } i < j \\ \bullet \quad \eta_{i}^{q+1}\eta_{j}^{q} = \eta_{j+1}^{q+1}\eta_{i}^{q} \text{ if } i \leq j \\ \bullet \quad \partial_{i}^{q+1}\eta_{j}^{q} = \eta_{j-1}^{q-1}\partial_{i}^{q} \text{ if } i < j \\ \bullet \quad \partial_{i}^{q+1}\eta_{i}^{q} = identity = \partial_{i+1}^{q+1}\eta_{i}^{q} \\ \bullet \quad \partial_{i}^{q+1}\eta_{j}^{q} = \eta_{j}^{q-1}\partial_{i-1}^{q} \text{ if } i > j+1 \\ \bullet \quad \lambda \in K^{q} \Rightarrow \eta_{i}^{q} \times \in K^{q+1} \\ \bullet \quad x \in K^{q} \Rightarrow \partial_{i}^{q} \times \in K^{q-1} \end{array}$$

Dependent parts

- face-gmsm
- invariant-gmsm

Independent parts

- degeneracy
- face-independent
- invariant-independent

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Functions

- face-absm
- degeneracy-absm
- invariant-absm
- Properties

$$\begin{array}{l} \textcircledleft{0.5} \partial_{i}^{q-1}\partial_{j}^{q}=\partial_{j-1}^{q-1}\partial_{i}^{q} \ \text{if} \ i < j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j+1}^{q+1}\eta_{i}^{q} \ \text{if} \ i \leq j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j-1}^{q-1}\partial_{i}^{q} \ \text{if} \ i < j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{i}^{q}=identity =\partial_{i+1}^{q+1}\eta_{i}^{q} \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j}^{q-1}\partial_{i-1}^{q} \ \text{if} \ i > j+1 \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j}^{q}\partial_{i+1}^{q} \ \text{if} \ i > j+1 \\ \textcircledleft{0.5} \chi \in K^{q} \Rightarrow \eta_{i}^{q} \chi \in K^{q-1} \\ \end{array}$$

Dependent parts

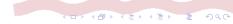
- face-gmsm
- invariant-gmsm

$$\ \, { \ \, 0 } \ \, \partial_i^{q-1} \partial_j^q { gmsm } = \partial_{j-1}^{q-1} \partial_i^q { gmsm }, i < j$$

• $gmsm \in K^q \Rightarrow \partial_i^q gmsm \in K^{q-1}$

Independent parts

- degeneracy
- face-independent
- invariant-independent
- Independent parts of Properties
 (1) and (7)



Functions

- face-absm
- degeneracy-absm
- invariant-absm
- Properties

$$\begin{array}{l} \textcircledleft{0.5} \partial_{j}^{q-1}\partial_{j}^{q}=\partial_{j-1}^{q-1}\partial_{i}^{q} \ \text{if } i < j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j+1}^{q+1}\eta_{i}^{q} \ \text{if } i \leq j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j-1}^{q-1}\partial_{i}^{q} \ \text{if } i < j \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{i}^{q}=identity = \partial_{i+1}^{q+1}\eta_{i}^{q} \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j}^{q-1}\partial_{i-1}^{q} \ \text{if } i > j+1 \\ \textcircledleft{0.5} \partial_{i}^{q+1}\eta_{j}^{q}=\eta_{j}^{q}\partial_{i-1}^{q} \\ \textcircledleft{0.5} X \in K^{q} \Rightarrow \eta_{i}^{q} \times \in K^{q+1} \\ \textcircledleft{0.5} X \in K^{q} \Rightarrow \partial_{i}^{q} \times \in K^{q-1} \end{array}$$

Dependent parts

- face-gmsm
- invariant-gmsm

$$\ \, { \ \, 0 } \ \, \partial_i^{q-1} \partial_j^q { gmsm } = \partial_{j-1}^{q-1} \partial_i^q { gmsm }, i < j$$

• $gmsm \in K^q \Rightarrow \partial_i^q gmsm \in K^{q-1}$

Independent parts

- degeneracy
- face-independent
- invariant-independent
- Independent parts of Properties
 (1) and (7)
- Properties (2) to (6)

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- Functions
 - face-absm
 - degeneracy-absm
 - invariant-absm
- Properties

$$\begin{array}{c|c} \bullet & \partial_{i}^{q-1}\partial_{j}^{q} = \partial_{j-1}^{q-1}\partial_{i}^{q} \text{ if } i < j \\ \bullet & \eta_{i}^{q+1}\eta_{j}^{q} = \eta_{j+1}^{q+1}\eta_{i}^{q} \text{ if } i \leq j \\ \bullet & \eta_{i}^{q+1}\eta_{j}^{q} = \eta_{j-1}^{q-1}\partial_{i}^{q} \text{ if } i < j \\ \bullet & \partial_{i}^{q+1}\eta_{i}^{q} = \eta_{i-1}^{q-1}\partial_{i-1}^{q} \text{ if } i > j \\ \bullet & \partial_{i}^{q+1}\eta_{i}^{q} = \eta_{j}^{q-1}\partial_{i-1}^{q} \text{ if } i > j + 1 \\ \bullet & \chi \in K^{q} \Rightarrow \eta_{i}^{q} \times \in K^{q+1} \\ \bullet & \chi \in K^{q} \Rightarrow \partial_{i}^{q} \times \in K^{q-1} \end{array}$$

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- Functions
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Concrete Simplicial Set

- 3 definitions + 7 theorems
- Proofs are not reusable for other cases

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Dependent parts

- face-gmsm
- invariant-gmsm
- $\partial_i^{q-1} \partial_i^q gmsm = \partial_{i-1}^{q-1} \partial_i^q gmsm, i < j$
- $gmsm \in K^q \Rightarrow \partial_i^q gmsm \in K^{q-1}$

Independent parts

- degeneracy
- face-independent
- invariant-independent
- Independent parts of Properties (1) and (7)
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Definitions of independent parts Proof of the independent theorems



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Definitions of independent parts Proof of the independent theorems Construction of a simplicial set instance



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Concrete Simplicial Set

- 2 definitions + 2 theorems
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(reduced) encapsulate + independent functions



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(reduced) encapsulate + independent functions

proof →

Generic Simplicial Set



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proof

Summary of our methodology

(reduced) encapsulate + independent functions ↑ encapsulate instance Generic Simplicial Set ↓ Concrete Simplicial Set + Proof correctness

- From 2 definitions and 2 theorems
- Instantiates 3 definitions and 7 theorems
- The proof of the 7 theorems involves: 92 definitions and 969 theorems
- The proof effort is considerably reduced

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- This methodology can be extrapolated to other cases

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Generic Theory for families of Simplicial Sets

A simplicial set

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A family of simplicial sets indexed by $\ensuremath{\mathcal{K}}$

(encapsulate

- ; Signatures
- (((imp-face-absm * * * *) => *)
- ((imp-degeneracy-absm * * * *) => *)
- ((imp-invariant-absm * * *) => *))
- ; Theorems
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A family of simplicial sets indexed by $\ensuremath{\mathcal{K}}$

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- ((inv-gmsm * * *) => *)

((indexp *) => *))

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(defthm imp-faceoface ;; $(\partial_i^{q-1} \partial_j^q gmsm = \partial_{j-1}^{q-1} \partial_i^q gmsm$ if i < j) ...)

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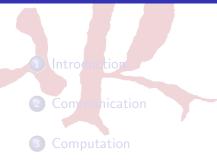
Generic Simplicial Set Theory

- Spheres family
- Standard Simplicial sets family
- Simplicial Complexes family

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4 Deduction

5 Conclusions and Further work



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Conclusions

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- Communication (*fKenzo*)
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- Deduction
 - Formalization of correctness of Kenzo programs using ACL2

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Further work

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 - Coq/SSReflect



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Mathematical Knowledge Management in Algebraic Topology

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May 31, 2011

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