

Formalising the structure of extremal p -groups

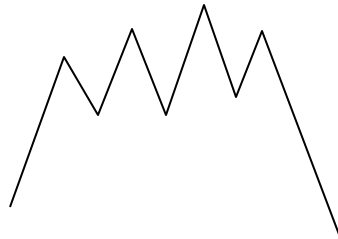
Georges Gonthier



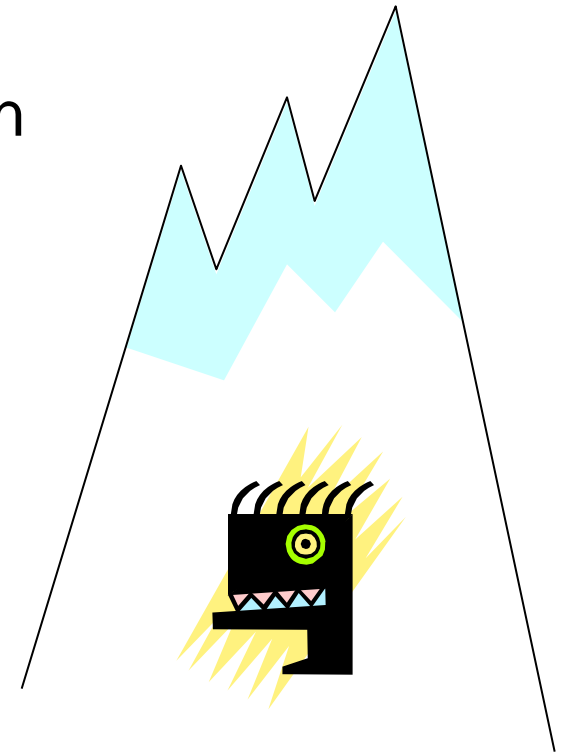
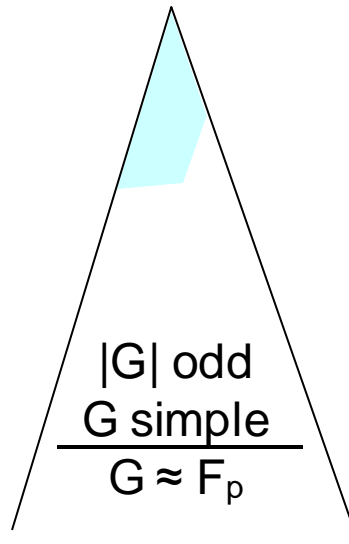
The Big Project

The Classification of Finite Simple Group

Frobenius groups
Thompson factorisation
character theory
linear representation
Galois theory
linear algebra
polynomials



Feit-Thompson



Goals and results

- Bring proof tools to mathematicians
- Apply software engineering to proof construction
- Prove landmark result
- Build on four-colour theorem experience
- Ssreflect proof language and Coq plugin
- Combinatorial components
- Linear algebra components
- Group theory components



4-color

Appel-Haken



finite groups

Feit-Thompson



Kepler

Hales



Tool review

- Data (inductive) types / propositions
- Computational reflection
 - compute values, types, and propositions
- Dependent types
 - first-class Structures
- Type / value inference
 - controlled by Coercion / Canonical Structure
- User notations



The Feit-Thompson proof

Let G be a minimal counter-example ...

Local Analysis

Character Theory

χ character: $\chi(g) = \text{tr}(\pi(g))$
for some $\pi : G \rightarrow M_n(\mathbb{C})$

Sylow: if $p^n \mid |G|$ then
 $|S| = p^n$ for some $S < G$

M_i maximal Frobenius
 $M_i \neq M_j$, K_i kernel of M_i

normal $S \triangleleft M$:
 $\forall m s, ms = s'm$ for some s'
 \hookrightarrow **factor group** M / S

$\in F_{p^q}[X]$
 $= q, P(F_p) = 0$
Galois
Theory

characters form a
euclidean space

kernel $K \triangleleft M$, $M = K \rtimes H$
if $hk = kh$ then $h = 1$

$|M_i| < |G|$
 $p < q, q < p$

impossible metrically



Local Analysis

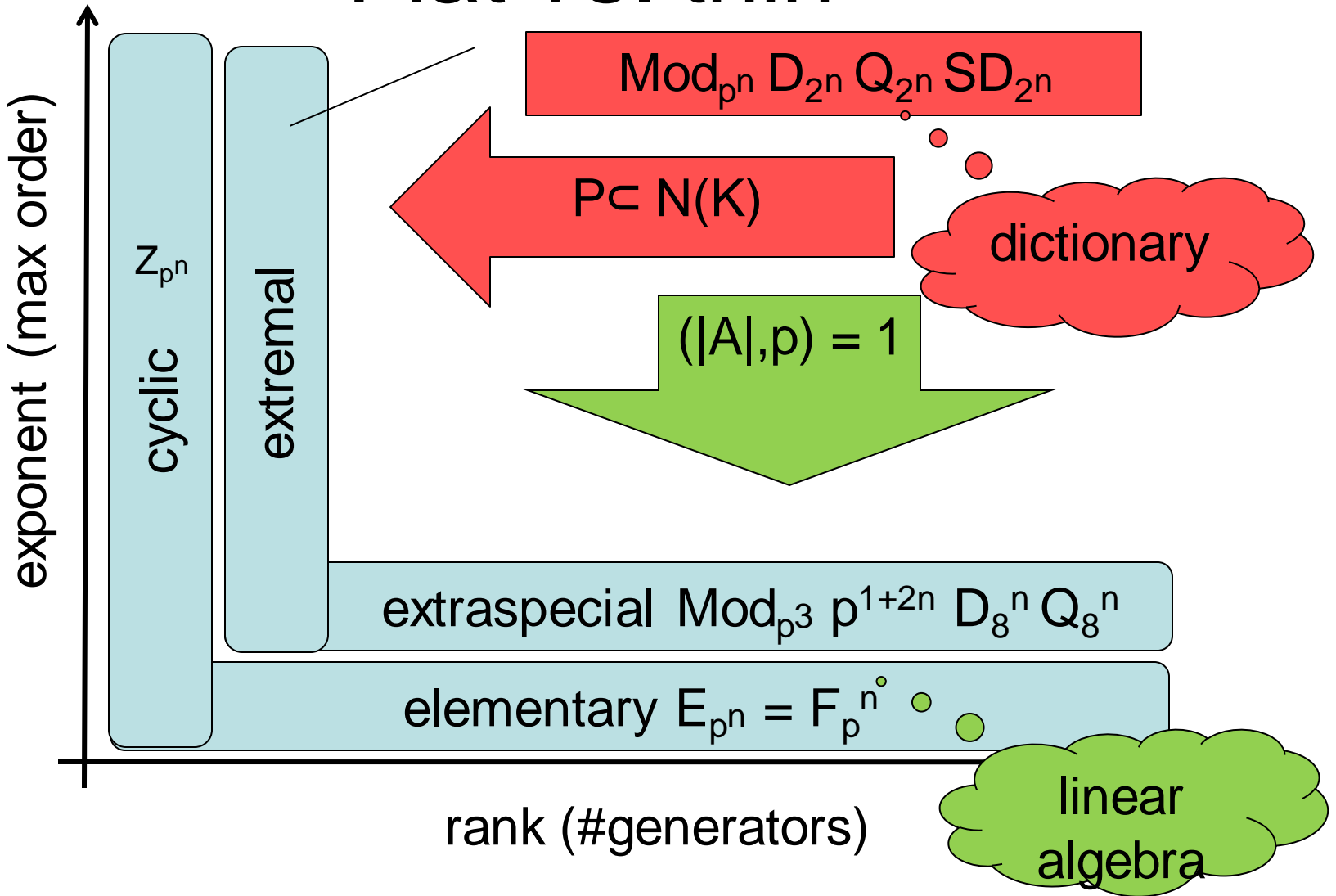
- Look at $N(P)$, P a p -group, i.e. $P \rtimes N(P)$
- Use the nilpotent structure of P do prove stuff by induction on P and/or $A \subset N(P)$
$$1 \triangleleft Z(P) \triangleleft Z_2(P) \triangleleft \dots \triangleleft P$$
- Base cases: extraspecial and **extremal p -groups**.

Roadmap

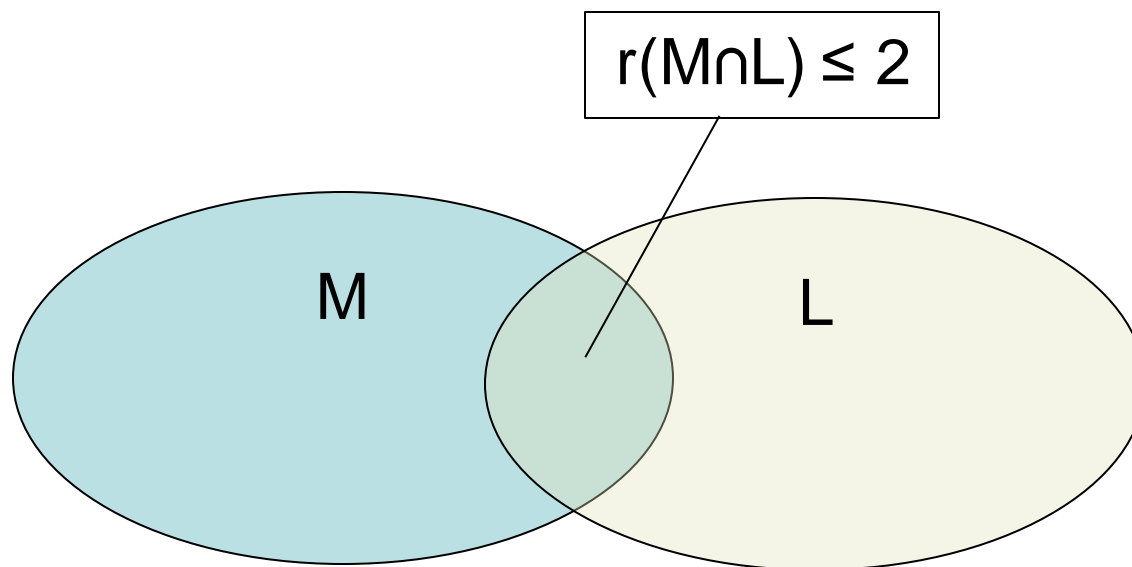
- Extremal p -groups
- Group presentations
- Big theorems



Flat vs. thin



The Uniqueness Theorem



Thompson transitivity:

$$S \in \text{Syl}_p(G), r(S) \geq 3$$

maximal abelian $A \triangleleft S$

Q_1, Q_2 maximal in $\mathcal{N}(A, q)$

$$Q_2 = Q_1^c \text{ for } c \in C_G(A)$$

p-stability:

$$\text{abelian } A \triangleleft S \in \text{Syl}_p(M)$$

for some $C \triangleleft M, p \nmid |C|$

$$A/C \triangleleft M/C$$

Dictionary

Theorem dihedral2_structure :

$n > 1 \rightarrow$ extremal_generators G 2 n $(x, y) \rightarrow G \setminus \text{isog } 'D_m \rightarrow$
 $[\wedge [\wedge X \times | Y = G, \{\text{in } G : \setminus : X, \text{forall } t, \#[t] = 2\}$
 $\& \{\text{in } X \& G : \setminus : X, \text{forall } z t, z^t = z^{-1}\}],$
 $[\wedge G^{(1)} = \langle [x^{+2}] \rangle, 'Phi(G) = G^{(1)}, \#|G^{(1)}| = r$
 $\& \text{nil_class } G = n.-1],$
 $'Ohm_1(G) = G \setminus (\text{forall } k, k > 0 \rightarrow 'Mho^k(G) = \langle [x^{+(2^k)}] \rangle),$
 $[\wedge yG : | : xyG = G : \setminus : X, [\text{disjoint } yG \& xyG]$
 $\& \text{forall } M, \text{maximal } M G = \text{pred3 } X \text{ My } Mxy M]$
 $\& \text{if } n == 2 \text{ then } (2.\text{-abelem } G : \text{Prop}) \text{ else}$
 $[\wedge 'Z(G) = \langle [x^{+r}] \rangle, \#|'Z(G)| = 2,$
 $\text{My } \setminus \text{isog } 'D_q, \text{Mxy } \setminus \text{isog } 'D_q$
 $\& \text{forall } U, \text{cyclic } U \rightarrow U \setminus \text{subset } G \rightarrow \#|G : U| = 2 \rightarrow U = X]].$

Presentations

- $D_{2^n} \approx \text{Grp} (x, y : x^{2^{n-1}}, y^2, xy^{-1})$
- $SD_{2^n} \approx \text{Grp} (x, y : x^{2^{n-1}}, y^2, xy^{-1})$
- $Q_{2^n} \approx \text{Grp} (x, y : x^{2^{n-1}}, y^2, xy^{-1})$

... but this is
not constructive (Post) !

Lemma Grp_dihedral :

'D_m \isog Grp (x : y : (x ^+ q, y ^+ 2, x ^ y = x ^-1)).

Faking free groups

- $G \approx \text{Grp } (x, y, \dots : r, s, t)$
→ $(H \simeq G) \Leftrightarrow (H \simeq \text{Grp } (x, y, \dots : r, s, t))$
- $H \simeq \text{Grp } (x, y, \dots : r, s, t)$
→ $(H = \langle x, y, \dots \rangle \ \& \ r = s = \dots = 1)$

Definition hom gT $(B : \{\text{set } gT\})$ $p :=$
 $\text{sat } B \ 1 \ \text{env1 } (p \ (Cst \ 0)).$

Definition iso gT $(B : \{\text{set } gT\})$ $p :=$
forall rT $(H : \{\text{group } rT\}), (H \setminus \text{homg } B) = \text{hom } H \ p.$

Faking relations

Inductive term :=

- | Cst of nat
- | Idx
- | Inv of term
- | Exp of term & nat
- | Mul of term & term
- | Conj of term & term
- | Comm of term & term.

Fixpoint eval {gT} e t : gT :=

match t with

- | Cst i => nth 1 e i
 - | Idx => 1
 - | Inv t1 => (eval e t1)^-1
 - | Exp t1 n => eval e t1 ^+ n
 - | Mul t1 t2 => eval e t1 * eval e t2
 - | Conj t1 t2 => eval e t1 ^ eval e t2
 - | Comm t1 t2 => [~ eval e t1, eval e t2]
- end.

Infix "*" := Mul : group_presentation.

Infix "^+" := Exp : group_presentation.

Infix "^" := Conj : group_presentation.

Notation "x ^-1" := (Inv x) : group_presentation.

Faking presentations

Inductive formula := Eq2 of term & term | And of formula & formula.

Inductive rel_type := NoRel | Rel vT of vT & vT.

Inductive type := Generator of term -> type | Formula of formula.

Inductive env gT := Env of {set gT} & seq gT.

Definition bool_of_rel r := if r is Rel vT v1 v2 then v1 == v2 else true.

Definition and_rel vT (v1 v2 : vT) r :=

if r is Rel wT w1 w2 then Rel (v1, w1) (v2, w2) else Rel v1 v2.

Fixpoint rel {gT} (e : seq gT) f r :=

match f with

| Eq2 s t => and_rel (eval e s) (eval e t) r

| And f1 f2 => rel e f1 (rel e f2 r)

end.

Using presentations

Lemma `Grp'_dihedral` :

`'D_m \isog Grp (x : y : (x ^+ 2, y ^+ 2, (x * y) ^+ q)).`

Proof.

`move=> gT G; rewrite Grp_dihedral.`

`apply/existsP/existsP=> [] [[x y]] /=.`

`case/eqP=> <- xq1 y2 xy; exists (x * y, y).`

`rewrite !xpair_eqE /= eqEsubset.`

...

`by rewrite y2 mulg1 xq1 !eqxx.`

`case/eqP=> <- x2 y2 xyq; exists (x * y, y).`

....

`by rewrite !mulgA mulgK -mulgA -(expG5 _ 1) x2 y2 mulg1.`

Qed.

Proving presentations

Lemma *isogEcard* :

forall rT aT (G : {group rT}) (H : {group aT}),
(G \isog H) = (G \homg H) && (#|H| <= #|G|).

Lemma *card_dihedral* : #|'D_m| = m.

Lemma *normal_rank1_structure* :

forall gT p (G : {group gT}),
p.-group G ->
(*forall* X : {group gT}, X <| G -> abelian X -> cyclic X) ->
cyclic G
∨ [&& p == 2, extremal2 G & (#|G| >= 16) || (G \isog 'Q_8)].

Big Theorems

Theorem dihedral2_structure:

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 $\& \{ \text{in } X \& G : \setminus : X, \text{forall } z t, z \wedge t = z \wedge^{-1} \}],$
 $[\wedge G \wedge (1) = \langle [x \wedge + 2] \rangle, 'Phi(G) = G \wedge (1), \# |G \wedge (1)| = r$
 $\& \text{nil_class } G = n.-1],$
 $'Ohm_1(G) = G / \setminus (\text{forall } k, k > 0 \rightarrow 'Mho \wedge k(G) = \langle [x \wedge + (2 \wedge k)] \rangle),$
 $[\wedge yG : | : xyG = G : \setminus : X, [\text{disjoint } yG \& xyG]$
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 $\text{My } \setminus \text{isog } 'D_q, \text{Mxy } \setminus \text{isog } 'D_q$
 $\& \text{forall } U, \text{cyclic } U \rightarrow U \setminus \text{subset } G \rightarrow \# |G : U| = 2 \rightarrow U = X]].$

Using big theorems

Lemma `generators_2dihedral`:

```
n > 1 -> G \isog 'D_m ->  
exists2 xy, extremal_generators G 2 n xy  
& let: (x, y) := xy in #[y] = 2 /\ x ^ y = x ^ -1.
```

Lemma `normal_rank1_structure`:

```
....  
have [[x y] genG _] := generators_2dihedral n_gt1 isoG.  
have [ _ _ [ _ _ maxG ] ] := dihedral2_structure n_gt1 genG isoG.
```

Conclusions

- Formalizing big math sometimes requires **big theorems**.
- It is paramount to build **usability** into formal theories.
- Type theory and reflection are powerful and flexible tools **engineer** this usability.



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